CSCI: 4500/6500 Programming Languages

Functional Programming Languages Part 4: Standard Meta Language (SML) & Haskell



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Standard ML

- ML historically stands for Meta Language. ML was a meta language for expressing and manipulating logical proofs.
- General purpose, modular functional programming language developed a team in the 1970s at the University of Edinburgh, headed by Robin Milner (polymorphism paper in reading list).
 - » Polymorphism one behavior for different types
 - » First language to include "polymorphic type inference" (functions with multiple different types – different input parameters) together with a type-safe exception handling mechanism.

Preview: Polymorphism

- One behavior (e.g., a function definition) for multiple different types (e.g., a function handles different types of input parameters).
- Ad-Hoc Polymorphism: range of actual types is finite and the combinations must be specified individually prior to use so multiple definitions (overloading, coercision) -> compiler calls the right definition
- Parametric Polymorphism (first type of polymorphism to appear in an actual programming language – ML in 1976)
 - » NO explicit type definition (e.g., the append function of a list)
 - » Used transparently with any number of types
 - » Generic Programming (arguably): Only one definition (e.g., templates, macro, note instantiation is laziness, "not evaluated until needed" characteristics)
- Suptyping Polymorphism (inheritance) classes related by Maria HySupertype.

Polymorphism

- Universal polymorphism : Allows writing code that works with different types
 - Ad-hoc polymorphism : Selecting the right implementation code to be executed



Standard ML

- Uses type declarations, but also does type inferencing to determine the types of undeclared variables
 » type of all variables can be determined at compile time.
 - » type of all variables can be determined
 » function Foo(a b) = a + b
- Static-scoped
- Syntax is closer to Pascal than to LISP
- » e.g., infix arithmetic expressions instead of Cambridge postfix
- Restrictions on how data types are intermixed (more later):
 » Example: integer division may not be used on strings
 - » ML is strongly typed (whereas Scheme is essentially typeless) and has no type coercions (talk more about this later in the Semester)
- Includes exception handling
- Module facility for implementing abstract data types.
- Permits side-effect (therefore an impure functional language)

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Standard ML (cont)

- Standard ML is a domain-specific language that is appropriate for building compilers
- Support for
 - » Complex data structures (abstract syntax, compiler intermediate forms)
 - » Memory management like Java
 - » Large projects with many modules
 - » Advanced type system for error detection

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Learn more details

- Today we will cover the basics so you can get started.
 Resources:
- » Robert Harper's (pdf book)
 - <u>http://www.cs.cmu.edu/~rwh/smlbook/online.pdf</u>
 » Peter Lee's:
 - <u>http://www.cs.cmu.edu/afs/cs/local/sml/common/smlguide/smlnj.htm</u>
 » SML/NJ Literature:
 - http://www.smlnj.org/doc/literature.html#tutorials
 - Runs on Microsoft Windows, MacOS X (yay!), UNIX,
 » Short & concise tutorial (not available)
 - http://cs.wwc.edu/Environment/SML-Tutorial.html
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Installation



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Distribution (SML of New Jersey):

- http://www.smlnj.org/dist/working/110.69/index.html
- Developed at Bell laboratories and Princeton University
- Installation (straight forward): Set your PATH variable where you install it.
 - » Forgot?
 - find / -name sml -print - /usr/local/smlnj-110.69/bin/sml # default on a MAC
 - /usr/iocal/sminj-110.69/bin/smi # default on a M - export PATH=\$PATH:/usr/local/smlnj-110.69/bin
 - export PATH \$PATH:/usi/iocal/smlnj=110.09/bin
 setenv PATH \$PATH:/usr/local/smlnj=110.69/bin

Run:

{saffron:ingrid:219} sml

Standard ML of New Jersey v110.69 [built: Tue Feb 3 22:24:07 2009]

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ML

Preliminaries

Interactive

- Type in expressions
- Evaluate and print type and result
- End with ;
- Exit (enter end of file) Compileable

• Read – Eval – Print – Loop:

Hello word! in SML

```
- print("Hello world!\n");
Hello world
val it = () : unit
-
```

"it" is the default name of the expression.

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Preliminaries

• Read - Eval - Print - Loop: - 3+2; val it = 5 : int

- 3+2;

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Preliminaries

• Read – Eval – Print – Loop:

- 3+2; val it = 5 : int

- it + 7 ;

Preliminaries

```
• Read - Eval - Print - Loop:
- 3+2;
val it = 5 : int
- it + 7 ;
val it = 12 : int
```

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Preliminaries

Read – Eval – Print – Loop:				
- 3+2;				
<pre>val it = 5 : int</pre>				
- it + 7 ;				
<pre>val it = 12 : int</pre>				
- it - 3 ;				

Preliminaries

```
• Read - Eval - Print - Loop:
- 3+2;
val it = 5 : int
- it + 7 ;
val it = 12 : int
- it - 3 ;
val it = 9 : int
```

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Preliminaries

```
• Read - Eval - Print - Loop:
- 3+2;
val it = 5 : int
- it + 7 ;
val it = 12 : int
- it - 3 ;
val it = 9 : int
- 4 + true
```

Preliminaries

```
• Read – Eval – Print – Loop:
   - 3+2;
   val it = 5 : int
   - it + 7 ;
   val it = 12 : int
    - it - 3 ;
   val it = 9 : int
   - 4 + true
    = ;
   stdIn:14.1-14.9 Error: operator and operand don't
      agree [literal]
     operator domain: int * int
                      int * bool
     operand:
     in expression:
4 + true
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```

• Copy and paste the following text into a Standard ML window:

2+2;	(* note semicolon at end*)
3*4;	
4/3;	(* an error! *)
6 div 2;	(* integer division *)
7 div 3;	

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```
- 4/3 ;
stdIn:20.1-20.4 Error: operator and operand
don't agree [literal]
operator domain: real * real
operand: int * int
in expression:
    4 / 3
- 4.0 / 3.0 ;
val it = 1.33333333333 : real
```

List functions

• [1,2,3,4];		Includes lists and list operations	
-val it : [1,2,3,4] : int list • val myList = [1,2,3,4]; -val myList : [1,2,3,4] : int list • 0 :: [1,2, 3]; -val it = [0,1,2,3] : int list - 213 :: 0 :: [1,2, 3];val it = [213,0,1,2,3] : int list		 The val statement binds a name to a value (similar to DEFINE in Scheme) Function declaration form: fun function_name (formal_parameters) = 	
		function_body_expression; e.a	
		fun cube(x : int) = $x * x * x$;	
		fun square(x: int) : int = $x * x$;	
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Haskell

- Similar to ML (syntax, static scoped, strongly typed, type inferencing)
- Different from ML (and most other functional languages) in that it is purely functional (e.g., no variables, no assignment statements, and no side effects of any kind)

Haskell

• Most Important features

- » Uses lazy evaluation (evaluate no subexpression until the value is needed)
- » Has list comprehensions, which allow it to deal with infinite lists

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Our First Program



Haskell

- Next project: you can choose Haskell or SML.
- For Haskell you are expected to use the Glasgow compiler (co-dependencies – perl, gcc) -> compiles like C (ghc –o main main.hs)
- Other compilers: Glasgow (Glorious), Helium, Hugs, Omega (is strict).

Concurrent Haskell

- ghc spare --make -threaded
- Enable threads:
 - » time ./primes-test +RTS -N2
- Run Example Program in threaded and unthreaded mode
 - » "True Parellelism"

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» Running threads in parallel & multiple processors – Pitfalls?

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Applications of Functional Languages

- APL is used for throw-away programs
- LISP is used for artificial intelligence
 - » Knowledge representation
 - » Machine learning
 - » Natural language processing
 - » Modeling of speech and vision
- Scheme is used to teach introductory programming at a significant number of universities

Comparing Functional and Imperative Languages

- Imperative Languages:
 - » Efficient execution
 - » Complex semantics
 - » Complex syntax
 - » Concurrency is programmer designed
- Functional Languages:
 - » Simple semantics
 - » Simple syntax

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- » Inefficient execution
- » Programs can automatically be made concurrent

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Function Definition

Functional Programming in Perspective (pros)

Functional Programming in Perspective (cons)

Advantages of functional languages

- » lack of side effects makes programs easier to understand
- » lack of explicit evaluation order (in some languages) offers possibility of parallel evaluation (e.g. MultiLisp)
- » lack of side effects and explicit evaluation order simplifies some things for a compiler (provided you don't blow it in other ways)
- » programs are often surprisingly short
- » language can be extremely small and yet powerful

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Advantages of functional languages

- » difficult (but not impossible!) to implement efficiently on von Neumann machines
 - lots of copying of data through parameters
 - (apparent) need to create a whole new array in order to change one element
 - heavy use of pointers (space/time and locality problem)
 - frequent procedure calls
 - heavy space use for recursion
 - requires garbage collection
 - requires a different mode of thinking by the programmer
 - difficult to integrate I/O into purely functional model

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