# COMP 110-001 <br> Recursion, Searching, and Selection 

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## Announcements

- Homework 4 deadline extended to June $13^{\text {th }}$, by $11: 59 \mathrm{pm}$
- Final exam, comprehensive
- Wednesday, June $17^{\text {th }}, 8 \mathrm{am}-11$ am
- Review on Monday


## Today

- Introduction to Recursion
- Introduction to Search \& Selection
- Not the focus of the final exam
- But, you should be able to understand the code in the slides (and know how to use the code in similar problems by making slight modifications).


## Recursion

- Whenever an algorithm has one subtask that is a smaller version of the entire algorithm's task, it is said to be recursive
- Recursion: you write a method to solve a big task, and the method invokes itself to solve a smaller subtask
- E.g., I want to eat 5 apples now. My subtask can be eating 4 apples, eating 3 apples, eating 2 apples, et.....
- To eat 5 apples, I can do:

Eat 1 apple

- Eat 3 apples + Eat 2 apples
- Eat 1 apple + Eat 4 apples

Eat 1 apple

## Recursion

- Eating 1 apple is the smallest task that I can have. I cannot divide it anymore.
- This is the base case in recursion.
- Recursion is to divide a big task into smaller tasks. Smaller tasks are then divided further. Until we reach base case.



## Recursion

- Let's start with a simple example: calculating factorial
- Factorial(n) $=\mathrm{n}$ * $(\mathrm{n}-1)$ * $(\mathrm{n}-2)$ * ..... * 3 * 2 * 1
- How do you solve a task with smaller task(s)?
- Factorial(n) = n * Factorial(n-1)



## Recursion

- Translate this into Java code public static int factorial( int n )
\{
if ( $n==1$ ) return 1; // base case
else



## Recursion

- The recursion form can be more natural in many problems (than using loops)
- Some problems can be hard to formulate using naïve looping (but such problems are beyond the scope of this course)
- Let's see more recursion examples:
- Digits to Words from textbook


## Recursion: Digits to Words

- Define a method that takes a single integer as an argument and displays the digits of that integer as words.
- For example, if the argument is the number 223, the method should display:
two two three
- Base case?
- Recursive rule?


## Recursion: Digits to Words

- Base case: only 1 digit
- print word for 1 digit
- Recursive rule:

Print words for n digits -->
(print words for first n-1 digits) + (print word for last digit)

## Recursion: Digits to Words

public static void displayAsWords( int number )
\{
if (number < 10) // base case
System.out.print(getWordFromDigit(number) + " ");
else //number has two or more digits
i
displayAsWords(number / 10);
System.out.print(getWordFromDigit(number \% 10) + " ");
\}
\}
You should be able to write out: getWordFromDigit(int num)

## Recursion: Digits to Words

```
\{//Code for invocation of displayAsWords (987)
    if (987 < 10)
            System.out.print(getWordFromDigit(987) + " ");
    else//987 has two or more digits
    \{
        _displayAsWords(987 / 10);
            System.out.print(getWordFromDigit(987 \% 10) + " ");
    \}
\}
```

- displayAsWords(987);

```
{//Code for invocation of disp7ayAsWords (98)
    if (98 < 10)
        System.out.print(getWordFromDigit(98) + " ");
    e7se//98 has two or more digits
    {
            displayAsWords(98 / 10);
                System.out.print(getWordFromDigit(98 % 10) + " ");
}
{//Code for invocation of displayAsWords(9)
    if (9 < 10)
        System.out.print(getWordFromDigit(9) + " ");
    e1se//9 has two or more digits
    {
        displayAsWords(9 / 10);
        System.out.print(getWordFromDigit(9 % 10) + " ");
    }
}
```


## Search

- Given a list of numbers (in an array), how do you search for a number?
- Return index if the number is found in the array
- Return -1 if the number is not found


## Sequential (Linear) Search

- Basic idea
- For each item in the list:
- if that item has the desired value, stop the search and return the item's location.
- Return Not Found.
- Can you do better than this (by making it faster)?
- The general answer is no
- No assumptions made on array (unsorted)
- In worst case, have to examine each array element at least once


## Search

- How about sorted array? ( numbers are in ascending or descending order )
- Can you make the linear search faster?


## Search

- Let's see an example: searching for 76



## Search

- Given n numbers:
- In linear search, you need to explore one possible choice in each iteration
- Worst case, n comparisons needed
- With the new search algorithm (which only works for sorted array), we can reduce half of the search space in each iteration!
- How many comparisons do I need in the worst case?



## Binary Search

```
int binary_search(int A[], int key, int imin, int imax) {
    // test if search range is empty
    if (imax < imin) {
    return KEY_NOT_FOUND; // set is empty
} else {
    // calculate midpoint to cut set in half
    int imid = midpoint(imin, imax);
    // three-way comparison
    if (A[imid] > key) // key is in lower subset
    return binary_search(A, key, imin, imid-1);
    else if (A[imid] < key) // key is in upper subset
            return binary_search(A, key, imid+1, imax);
    else // key has been found
    return imid;
    }

\section*{Search Algorithms}
- A lot of search algorithms, here we just covered two simplest cases:
- Linear search in a list (array) of numbers
- Binary search in sorted array
- More with different data structures:
- Search in graphs and trees (computer science concepts, not the usual graph/tree)
- E.g., search for a move in chess game
- Search for relations/patterns in social network communication graph

\section*{Selection}
- One selection problem:
- Find the smallest / largest number in a given list (array)
- No assumption made on the list ( so it is not sorted)
- We have solved this in lab 4
- Loop through each element, keep the largest/ smallest
- Let's relax the problem a bit

\section*{Selection}
- Find the k-th smallest (or largest) element in a list of numbers
- How to solve this problem?
- Go through each element, for each element, check its position in list
- How many operations in the worst case?
- Sort array first. Then get the \(k\)-th element
- How many operations in the worst case

\section*{Selection}

Quickselect (quick in practice, but not in the worst case)
- To find \(\mathbf{k}\)-th smallest number in n numbers:
- Randomly pick a number from the list, call it \(p\)
- Partition the array into two parts:
- Numbers that are < p (m numbers)
- Numbers that are >p ( \(n-m-1\) numbers)
- If \(m==k-1, p\) is the \(k\)-th smallest
- If \(m>k\), find the \(k\)-th smallest in the \(m\) numbers
- If \(m<k\), find the ( \(k-m-1\) )-th smallest in the ( \(n-m-1\) ) numbers
- On average, this requires \(\sim n^{*}\) constant operations
- But in the worst case, it is \(\sim n^{\wedge} 2^{*}\) constant

Next Class
- Introduction to sorting```

